Zero-Knowledge Proofs Showing you *can* do something, without showing how

Alex Dehnert

6.UAT

Fall 2011

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Zero-Knowledge Proofs

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Problem: Ali Baba's Cave



- Imagine that you (Peggy) are exploring a cave with a friend
- You know a secret password to get through the door between R and S
- You want to prove to your friend Victor that you know the password, without him finding out the password.

Image: RSA Laboratories

Solution

- Victor stands outside at P
- Peggy vanishes into the cave and stands at R or S
- Victor comes back in, stands at Q, and calls out which side Peggy should return by
- Peggy returns as requested



What if Peggy cheats?

- Victor stands outside at P
- Peggy picks R or S to hide at
- Victor comes back in, stands at Q, and calls out which side Peggy should return by
- If Peggy picked the correct side, she tricks Victor; otherwise, she has to admit she lied
- If Victor does this enough times, Peggy will guess wrong (one trial, 50% she succeeds; two trials, 25%; etc.)



- Goal: Peggy (prover) wants to convince Victor (verifier) of something Convincing If true, an honest prover can always convince an honest verifier
- No cheating If false, a dishonest prover can only rarely convince an honest verifier
- Zero knowledge If true, a cheating verifier cannot learn anything more than that it is true.

Peggy: Prepare

Peggy: Commit

Victor: Challenge

Peggy: Respond

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Victor	stands	outside
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Peggy: Prepare

Peggy vanishes into the cave

Peggy: Commit

Victor: Challenge

Peggy: Respond

	Victor stands outside
Peggy: Prepare	Peggy vanishes into the cave
Peggy: Commit	Victor comes back in,
Victor: Challenge	

Peggy: Respond

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	Victor stands outside
Peggy: Prepare	Peggy vanishes into the cave
Peggy: Commit	Victor comes back in,
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Peggy: Respond

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	Victor stands outside
Peggy: Prepare	Peggy vanishes into the cave
Peggy: Commit	Victor comes back in,
Victor: Challenge	and calls out which side Peggy should return by
Peggy: Respond	Peggy returns as requested

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Map 3-coloring

Given a map, prove that it can be colored using only three colors (and that you know a coloring).





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- Numerous cases where you want to prove who you are
- Example: authorize a credit card transaction
- Don't want stores to be able to charge you for more stuff later
- Use zero-knowledge proof to prove who you are, instead of a credit card number you keep secret

Using VProbes for Intrusion Detection

Alex Dehnert

VMware Intern, Summer 2011 Massachussetts Institute of Technology, S.B. 2012

Thursday, September 15, 2011

Virtualization



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Virtualization



Hardware

- Security
- Unique virtualization features
 - vprobes

vprobes IDS



Use vprobes to watch application/OS communication

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- Build a profile of normal behavior
- Alert when abnormal behavior occurs

Cost of Intrusions

\$114 billion Total global cost of cybercrime annually, as estimated by Symantec

- \$75 billion 2011 spending on IT security for US companies, according to the Ponemon Institute
- \$170 million Cost to Sony of the recent PlayStation Network hack, including shutting down PSN, network security improvements, etc.

Sources: The Fiscal Times, Huffington Post

Intrusion Detection in Virtual Machines

Alex Dehnert

VI-A - VMware

Fall 2011

- Security is a growing issue
- Huge amounts of money are being spent, both to defend against attacks and to clean up after them
- Host-based intrusion detection systems monitor servers to detect attacks
- Security software is sometimes specifically targeted

System calls (syscalls)



Intrusion detection

- Rich variety of work involving host-based intrusion detection
- Many based on patterns of system calls •



Applications

How can virtual machines help?



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Provide an intrusion detection system that is robust against attacks from malware, by leveraging vprobes to run the agent on the host

Much previous work on how to determine whether or not some sequence of syscalls is allowed

- Hand-written to match whatever the server in question is doing
- "Only allow the web server to access files under /var/www/"
- "Only allow the ssh server binary to bind to port 22"

- Maintain current "state" associated with the syscalls so far (and other info)
- Define transitions between the various states
- When the appropriate syscall is seen, transition to the appropriate new state:



• If an unfamiliar syscall is made, alert

- Sequence Time-Delay Embedding
- Sequences (*n*-grams) of syscalls



• Look at rolling window of twenty sequences at a time. If enough of them are unknown, alert.

- Intrusion detection based on syscalls *isn't* plenty of people have done that
- However, generally, that's vulnerable to attacks against the agent
- So, develop a version that gathers data using vprobes

- vprobe script to get syscall information from the kernel (Summer 2011)
- Build analysis modules (begun Summer 2011)
- Build up collection of exploits and normal behaviors to test against

Timeline

- Week 1 Get set up at VMware again
- Week 2 Re-familiarize with vprobes
- Week 3 Find off-the-shelf exploits for Apache and other services
- Week 4 Write a custom Apache exploit
- Week 5 More exploit-writing
- Week 6 Review literature on intrusion detection again; begin writing analyzer (e.g., finite automata)
- Week 7 Continue literature-based analyzer
- Week 8 Start custom analyzer (e.g., SVM)
- Week 9 Continue custom analyzer
- Week 10 Code cleanup and begin code review
- Week 11 Finish code review; begin writing report
- Week 12 Finish writing report

Big risk: schedule slippage Room to cut in both the exploit finding/writing and the analyzers Can't come up with a novel analyzer Not key to the project. The main novelty is supposed to be in the vprobes component, anyway. vprobes not powerful enough My team is the vprobes team. We can probably add required features.

- Security becoming increasingly important
- Host-based intrusion detection vulnerable to the agent being attacked
- Project: Enhance effectiveness of intrusion detection systems by moving the vulnerable agent onto the host, and use vprobes to look inside the VM
- Extensive literature exists on how to examine the data the agent/vprobes gather
Compton Scattering and the Klein-Nishina Formula

Daniel J. Fremont

8.13 Massachusetts Institute of Technology

History

- Compton's x-ray diffraction experiments (1920-2)
 - Wavelength shift in scattered radiation
 - Explanation using photon momentum
- Wilson, Bothe detect recoil electron (1923)
- Wave-particle duality (de Broglie, Schrödinger)

Compton Scattering Theory

- Billiard-ball style collisions
- Relativistic energy: $E^2 = p^2 c^2 + m^2 c^4$



Scattering Intensities

- Classical treatment: Thomson cross-section
 - Only valid for low energies
- Relativistic treatment: Klein-Nishina formula



Apparatus



- Discriminators filter out noise
- Coincidence detector selects Compton events

Data



Amplitudes



Amplitude Results

• Classically-expected increase for large θ not observed

- Large systematic errors
 - Unusually weak signal at 30° likely discriminated against
 - Scintillator efficiency
 - Detector drift
- Errors too significant to verify Klein-Nishina; however data definitely inconsistent with Thomson model

Attenuation



Attenuation Results

• Measured electron cross-section: $(2.50 \pm 0.025) \cdot 10^{-29} \, m^2$

- Klein-Nishina prediction: $2.56 \cdot 10^{-29} m^2 (\approx 2 \sigma \text{ off})$
- Thomson prediction: $6.65 \cdot 10^{-29} m^2 (\approx 166 \sigma \text{ off!})$
- Small systematic error
 - Some scattered photons detected
 - Detector drift
- Relativistic effects dominant: Klein-Nishina model much more accurate than classical Thomson model

Summary

- Systematic errors in amplitude data too large to allow verification of Klein-Nishina, but data definitely inconsistent with classical prediction
- Attenuation data strongly supports Klein-Nishina result over Thomson
- Necessity of relativistic quantum description of photons and electrons

Questions?

- History of Compton Scattering
- Scattering Theory
- Amplitude Experiment
- Attenuation Experiment

Electron–Xenon Scattering: The Ramsauer-Townsend Effect

Daniel J. Fremont

8.13 Massachusetts Institute of Technology

History of Ramsauer-Townsend

- Ramsauer (1921) discovers unexpected behavior in scattering experiments
- Townsend and Bailey (1922) confirm effect using a different method
- Bohr proposes interference of matter waves
- Effect explained using Schrödinger equation (1927-31)

Scattering



• $d\sigma = D(\theta) d\Omega$ • $\sigma(E) = \int D(\theta) d\Omega$

Matter Waves and Interference

- Electrons and atoms cannot be treated as billiard balls
- Objects are diffuse, described by a wave function
- Interference between wave functions can produce unexpected effects
- Ramsauer-Townsend: $\sigma = 0$ for a particular energy

Apparatus



Can freeze out Xe using liquid nitrogen

Fremont (MIT)

Ramsauer-Townsend

Calculating the Cross-Section

• Uniform gas density \rightarrow constant scattering probability for a fixed distance

$$I = I_0 e^{-\frac{L}{\lambda}}$$
$$\frac{1}{\lambda} = n\sigma$$

$$\sigma(V) = \frac{1}{nL} \ln \left[\frac{I_0(V)}{I(V)} \right]$$

Cross-Section Measurements



Results

- Minimum σ occurs at 0.53 \pm 0.077 V
 - Significantly off from true value of 0.9 V
- Random errors insufficient
- Systematic errors
 - Contact potential
 - Nonzero initial electron energy
 - Unequal filament temperatures
- Corrected value: 0.88 \pm 0.092 V

Improvements and Extensions

- Additional experiments to determine contact potential and initial electron energy
- Collect more data in vicinity of minimum σ
- Determine voltage necessary to get equal filament temperatures

Summary

- Observed the Ramsauer-Townsend effect in electron–Xenon collisions
- Found minimum cross-section near expected energy
- Provided evidence for interference in matter waves and quantum theory

Questions?

- History of Ramsauer-Townsend
- Scattering
- Matter Waves
- Dependence of cross section on energy

Richardson's Law

- $I \propto T^2 e^{-W \over kT}$
- Only radiative heat transfer ightarrow $V^2 \propto$ T^4



Child's Law • $I \propto V^{\frac{3}{2}}$



Nuclear Magnetic Resonance: Effect of Paramagnetic Ions on Relaxation Times

Daniel J. Fremont

8.13 — Junior Lab Massachusetts Institute of Technology

Outline

- Nuclear Magnetic Resonance
- Relaxation and Paramagnetic lons
- A Useful Tool

- Rabi (1937) predicts and first observes NMR
- Purcell and Bloch (1945) measure relaxation times and their dependence on chemical properties
- Chemical shifts discovered (1950-1)
- Lauterbur (1973) applies NMR to medical imaging

Pulsed NMR

- Nuclear spins precess around static magnetic field
- Oscillatory perpendicular field → nutation
- Perpendicular spin component induces measurable current



Spin-Lattice Relaxation (T1)

- Return of spins to thermal equilibrium
- Energy transferred into molecular motions
- Lattice frequencies must be close to Larmor frequency

Spin-Spin Relaxation (T2)

- Loss of phase coherence
- Interactions between nearby spins

Paramagnetic lons

- Enormous magnetic moments (~1000 times proton)
- Much larger interaction energies \rightarrow smaller relaxation times
- $\bullet\,$ Slow local field variations \to additional T2 mechanism

- $\frac{1}{T_1} \propto \text{ion concentration}$
- $\frac{1}{T^2}$ more strongly affected



Spin Echoes

Measuring relaxation directly has problems

- Proximity to driving pulse
- Inhomogeneity in static field alters observed T2
- Use 180° pulses to produce spin echoes (Hahn 1950)
 - Reverses phase decoherence from inhomogeneity



Apparatus



Finding Resonance

- Minimize beats
 - $\nu_{beats} = |\nu_{pulses} \nu_{Larmor}|$





- Find most uniform B field
 - Inhomogeneity \rightarrow different resonances \rightarrow weaker signal
- Result: $(1.424 \pm 0.083) \cdot 10^{-26}$ J/T (3% off accepted value)

T1 Measurement



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T2 Measurement



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Effect of Fe^{3+} lons on T1 and T2



Results

- Proton magnetic moment: $(1.424 \pm 0.083) \cdot 10^{-26} \text{ J/T}$
 - Accepted value: 1.411 · 10⁻²⁶ J/T
- $\frac{1}{71}$ grows linearly with Fe^{3+} concentration
- $\frac{1}{T^2}$ grows faster than linearly with Fe^{3+} concentration

Improvements and Extensions

- More stable probe setup to allow more accurate determinations of field strength
- Greater sensitivity to allow use of higher concentrations
- Removing dissolved oxygen



- Observed NMR in protons at the expected frequency
- Verified expected effects of *Fe*³⁺ ions on relaxation times
- Suggested usefulness of NMR in chemical analysis

Questions?

- History of NMR
- Pulsed NMR
- Spin-Lattice and Spin-Spin Relaxation
- Spin Echoes
- Proton Larmor Frequency
- Effects of *Fe*³⁺ ions on Relaxation Times

Paramagnetism: N₂ vs. O₂



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NMR - Ions and Relaxation Times

Observation of the Photoelectric Effect

Daniel J. Fremont

Massachusetts Institute of Technology

Fremont	(MIT	

History of a phenomenon

Discovery

- Hertz's radio experiments (1886-7)
- Ultraviolet light and electricity

Investigation

- Lenard's experiment (1902)
- Energies independent of intensity
- But dependent on frequency!

Understanding

- Einstein's explanation (1905)
- Light quanta with $E \propto \nu$
- Millikan's experiments (1912-5)

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Einstein's theory

- Planck's quantized oscillators
- Photons, with $E = h \nu$
- Intensity \equiv number of photons
- One photon, one electron

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- Energy of emitted electrons independent of intensity
- Energy dependent on frequency
- Minimum cutoff frequency

Einstein's theory

$$KE_{e^-} = h\nu - \phi$$



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Experiment Setup



Slightly modified from experiment lab guide (http://web.mit.edu/8.13/www/intro2.shtml).

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Current vs. Retarding Voltage



- Some expected features, with some extras
- Theoretical model?

Cutoff Voltage vs. Frequency



- Large χ^2 for small ν
- However, near-linear relationship as expected

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Results and Errors

- $h = (6.348 \pm 0.0626) \cdot 10^{-34} \text{ J} \cdot \text{s}$
 - \approx 4.2% off of true value (\approx 4.5 times estimated error)
- $\phi = 1.76 \pm 0.026 \text{ eV}$
 - $\blacktriangleright~\approx$ 23% off of true value (\approx 21 times estimated error)
- Random Errors
 - Precision of instruments
 - Small compared to...
- Systematic Errors
 - Ambient light
 - Light striking the anode
 - Cutoff voltage bias
 - Variation of ϕ
- Various possibilities for systematics, but probably not due to measurement error
- Disparity in error makes some possibilities less likely

Summary

- $h = (6.348 \pm 0.0626) \cdot 10^{-34} \text{ J} \cdot \text{s}$
- $\phi = 1.76 \pm 0.026 \text{ eV}$
- Photoelectric effect well described by Einstein's model
 - Linear dependence of cutoff voltage on frequency
- Particle-like nature of light

Questions?

Rock Paper Scissors:

Real Professional Strategy Nathan Benjamin & Xavier Jackson (nathanb@mit.edu, jaxxson@mit.edu)

Rock







Scissors



Avalanche (RRR)



Bureaucrat (PPP)



Edward's Gambit (SSS)





Crescendo (PSR), Decrescendo (RSP)



Scrapbook (PSP)



Fistful of Dollars (RPP)



Sally Wolfe

July 30, 2010

- Introduction

abc-permutations

abc-permutations

Definition

Let *abc*-permutations be elements of S_n obtained by partitioning [n] into three blocks of length a, b, c and exchanging the first and last blocks.

- Introduction

abc-permutations

abc-permutations

Definition

Let *abc*-permutations be elements of S_n obtained by partitioning [n] into three blocks of length a, b, c and exchanging the first and last blocks.

Examples		
[67851234] [67812345]		

-Introduction

abc-permutations

Work of Pak and Redlich

In 2008, Pak and Redlich investigated the probability that an *abc*-permutation of length n is a single n-cycle in order to answer a question posed by V. I. Arnold.

- Introduction

abc-permutations

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Theorem (Pak-Redlich, 2008)

Let $\sigma_{\rm abc}$ be an abc-permutation. Then $\sigma_{\rm abc}$ is a long cycle if and only if

gcd(a+b, b+c) = 1.

Introduction

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$$gcd(a+b, b+c) = 1.$$

Theorem (Pak-Redlich, 2008)

Let $\mathbf{p}(n)$ be the probability that an abc-permutation of length n is a long cycle. Then

$$\lim_{\to\infty}\mathbf{p}(n)=\frac{6}{\pi^2}.$$

n

Results on cycle structure of *abc*-permutations

Cycle structure of *abc*-permutations

We fully characterize the cycle structure of *abc*-permutations.

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Results on cycle structure of *abc*-permutations

Cycle structure of *abc*-permutations

We fully characterize the cycle structure of *abc*-permutations.

Theorem

Let k = gcd(a + b, b + c). Then σ_{abc} is composed of k cycles of lengths $\lfloor \frac{n}{k} \rfloor$ and $\lceil \frac{n}{k} \rceil$. These cycles correspond to the residue classes of n modulo k.

Results on cycle structure of *abc*-permutations

Cycle structure statistics of *abc*-permutations

We investigate the probability $\mathbf{p}_k(n)$ that a random *abc*-permutation of length *n* has exactly *k* cycles.

Theorem

We have

$$\lim_{n\to\infty}\mathbf{p}_k(n)=\frac{1}{k^2}\frac{6}{\pi^2}.$$

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Results on cycle structure of *abc*-permutations

Sketch of cycle structure classification

Ideas from the classification of cycle structure

First, note that if $x \leq c$, we have

$$\sigma_{abc}(x) = x + a + b.$$

For ease of notation, let d = c - a. If $c < x \le b + c$, we have

$$\sigma_{abc}(x) = x - d,$$

and if x > b + c, we have

$$\sigma_{abc}(x) = x - a - b.$$
Results on cycle structure of *abc*-permutations

Sketch of cycle structure classification

Ideas from the classification of cycle structure, cont

Therefore, for all i and x we have

$$\sigma^i_{abc}(x) = x + m(a+b) - ld.$$

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for some integers m, I satisfying $|m| + |I| \le i$.

Results on cycle structure of *abc*-permutations

Sketch of cycle structure classification

Ideas from the classification of cycle structure, cont

Therefore, for all i and x we have

$$\sigma^i_{abc}(x) = x + m(a+b) - ld.$$

for some integers m, I satisfying $|m| + |I| \le i$.

Let k = gcd(a + b, d) = gcd(a + b, b + c). Then the orbit of x under σ_{abc} will stay within one residue class modulo k.

Results on cycle structure of *abc*-permutations

Sketch of cycle structure classification

Ideas from the classification of cycle structure, cont

To show that each residue class contains exactly one residue class, we note that if x is in a cycle of length j, then

$$\sigma^{j}_{abc}(x) = x + m(a+b) - ld = x.$$

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and $m + l \leq j$.

Results on cycle structure of *abc*-permutations

Sketch of cycle structure classification

Ideas from the classification of cycle structure, cont

To show that each residue class contains exactly one residue class, we note that if x is in a cycle of length j, then

$$\sigma^{j}_{abc}(x) = x + m(a+b) - ld = x.$$

and $m + l \le j$. We show that $\lceil \frac{n}{k} \rceil / 2 < j$.

Results on cycle structure of *abc*-permutations

Cycle structure statistics

Idea of the construction of $\mathbf{p}_k(n)$

Let A_k be the event that k|(a + b) and (b + c), and let B_k be the event that (a + b)/k and (b + c)/k are relatively prime integers. Then

$$\mathbf{p}_k(n) = Pr(A_k) \cdot Pr(B_k \text{ given } A_k)$$

Results on cycle structure of *abc*-permutations

Cycle structure statistics

Idea of the construction of $\mathbf{p}_k(n)$, cont

Let f(n, k) be the probability that $k \mid (a + b)$ and (b + c) for a random *abc*-permutation of length *n*.

Lemma (Pak-Redlich, 2008)

$$f(n,k) = \begin{cases} \frac{\binom{n}{k}+1\binom{n}{k}+2}{(n+1)(n+2)} & \text{if } k \mid n \\ \frac{\lfloor \frac{n}{k} \rfloor \lfloor \lfloor \frac{k}{k} \rfloor + 1 \end{pmatrix}}{n(n+1)} & \text{if } k \nmid n \end{cases}$$

Results on cycle structure of *abc*-permutations

└─ Cycle structure statistics

Idea of the construction of $\mathbf{p}_k(n)$, cont

Then

$$\mathbf{p}_k(n) = f(n,k) \cdot Pr(B_k \text{ given } A_k).$$

Results on cycle structure of *abc*-permutations

Cycle structure statistics

Idea of the construction of $\mathbf{p}_k(n)$, cont

Then

$$\mathbf{p}_k(n) = f(n,k) \cdot Pr(B_k \text{ given } A_k).$$

If $n \mid k$, we have that $k \mid (a + b)$ and (b + c) implies that $k \mid a, b, c$. Then choosing the blocks a, b, c partitioning n is equivalent to choosing blocks a/k, b/k, c/k partitioning n/k.

Results on cycle structure of *abc*-permutations

Cycle structure statistics

Idea of the construction of $\mathbf{p}_k(n)$, cont

Then

$$\mathbf{p}_k(n) = f(n,k) \cdot Pr(B_k \text{ given } A_k).$$

If $n \mid k$, we have that $k \mid (a + b)$ and (b + c) implies that $k \mid a, b, c$. Then choosing the blocks a, b, c partitioning n is equivalent to choosing blocks a/k, b/k, c/k partitioning n/k.

Then $Pr(B_k \text{ given } A_k) = \mathbf{p}(n/k)$, so

$$\mathbf{p}_k(n) = f(n,k)\mathbf{p}(n/k).$$

- Results on cycle structure of *abc*-permutations
 - └─ Cycle structure statistics

Definition of $\mathbf{p}_k(n)$

Since

$$\mathbf{p}(n) = 1 - \sum_{s=2}^{n} \mu(s) f(n,s),$$

we have

- Results on cycle structure of *abc*-permutations
 - └─ Cycle structure statistics

Definition of $\mathbf{p}_k(n)$

Since

$$\mathbf{p}(n) = 1 - \sum_{s=2}^{n} \mu(s) f(n,s),$$

we have

Theorem

$$\mathbf{p}_{k}(n) = \begin{cases} f(n,k)(1 - \sum_{s=2}^{n/k} \mu(s)f(n/k,s)) & \text{if } n \mid k \\ f(n,k)(1 - \sum_{s=2}^{n/k} \mu(s)\Gamma(n/k,s)) & \text{if } n \nmid k \end{cases}$$

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where $\Gamma(n, k)$ is a similar function to f(n, k).

- Results on cycle structure of *abc*-permutations
 - Cycle structure statistics

Limit of
$$\mathbf{p}_k(n)$$

Theorem

We have

$$\lim_{n\to\infty}\mathbf{p}_k(n)=\frac{1}{k^2}\frac{6}{\pi^2}.$$

Pattern avoidance and *abc*-permutations

Pattern avoidance

Let
$$\sigma \in S_n$$
, and $\phi \in S_k$ where $k \leq n$.

Definition

We say that σ contains ϕ if there exists a subsequence of σ which is order isomorphic to ϕ . Otherwise, we say that σ avoids ϕ .

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Pattern avoidance and *abc*-permutations

Pattern avoidance and *abc*-permutations

Theorem

The set of abc-permutations is the set of permutations which avoid 132, 213, and 4321.

Pattern avoidance and *abc*-permutations

Pattern avoidance and *abc*-permutations

Theorem

The set of abc-permutations is the set of permutations which avoid 132, 213, and 4321.

In order to prove this theorem we make use of the fact that *abc*-permutations are *reverse layered permutations* with three or fewer layers.

Pattern avoidance and *abc*-permutations

Reverse layered permutations

Definition

A permutation ϕ is called reverse layered if it is of the form $q_1q_2 \dots q_k$, where the q_i are strings of consecutively increasing numbers and $q_i > q_j$ if i < j.

Pattern avoidance and *abc*-permutations

Reverse layered permutations

Definition

A permutation ϕ is called reverse layered if it is of the form $q_1q_2 \dots q_k$, where the q_i are strings of consecutively increasing numbers and $q_i > q_j$ if i < j.

Lemma (Monsour, 2002)

The set of reverse layered permutations is the set of 132 and 213 avoiding permutations.

Pattern avoidance and *abc*-permutations

Proof of classification of *abc*-permutations

Assume that σ is reverse layered and avoids 4321. Then it has at most three layers, so it is an *abc*-permutation.

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Pattern avoidance and *abc*-permutations

Proof of classification of *abc*-permutations

Assume that σ is reverse layered and avoids 4321. Then it has at most three layers, so it is an *abc*-permutation.

Now assume that σ is an *abc*-permutation. Then it is reverse layered, so it avoids 132, and 213. It also has at most three layers, so it avoids 4321.

Pattern avoidance and *abc*-permutations

Avenues for further research

1 What is the cycle structure of layered permutations?

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Pattern avoidance and *abc*-permutations

Avenues for further research

- 1 What is the cycle structure of layered permutations?
- 2 What is the structure and number of permutations which avoid 132, 213, and another permutation of length at least four?

Stabilizing unstable systems: balance and levitation

Sally Wolfe



Filling a glass to a specific height



Filling a glass to a specific height



Constant rate of flow

Time = volume/rate of flow

Filling a glass to a specific height



Choosing behavior based on error

- Desired result
- Method of making error measurement
- Way to control the system
- Way of deciding what to do based on error



Unstable systems



Unstable systems



Levitation!





Position sensor

High Speed Crystallization



Crystallization of supersaturated Sodium Acetate solution from below. The two points of emanation caused by two seeds are clearly visible. Shot at 1000 fps, f/8, 90mm lens.

Robin Deits, Sally Wolfe, Daniel Ron, Bio Lili

11/8/2010 • 1

Background – Goals – Procedure – Results – Discussion

Background: crystallization

- Crystallization refers to the formation of solid crystals from a homogeneous solution.
- It is essentially a solid-liquid separation technique and a very important one at that.
- Solubility Hot liquid dissolves more compounds. Once cooling process starts, compounds become crystals

Background – Goals – Procedure – Results – Discussion

Background: sodium acetate

- Sodium acetate (NaAc) can perform rapid crystallization when supersaturated
- Supersaturation is when more solute than is normally possible is dissolved into a solvent
 - Can be achieved by increasing temperature
- If a supersaturated solution of NaAc is slowly cooled, it can be rapidly crystallized if perturbed or given a nucleation site



[1] http://jchemed.chem.wisc.edu/JCESoft/CCA/CCA3/MAIN/ACETATE/PAGE1.HTM

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Background – Goals – Procedure – Results – Discussion

Background: bismuth

- White/silverish, crystalline, brittle, very dense, highly diamagnetic metallic element
- Very sensitive to high temperatures
- Very viscous and cools rapidly
- Pretty colors and cool shapes





1. <u>http://www.facts-about.org.uk/science-element-bismuth.htm</u>

2. <u>http://upload.wikimedia.org/wikipedia/commons/6/65/Bismuth_Crystal.jpg</u>

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Goals:

- Observe crystallization process
 - Targets: Sodium Acetate, Bismuth
- Make measurements of speed and process of crystallization
- Determine feasibility of future quantitative studies
Procedure: underneath setup



Procedure: underneath setup



Light

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Procedure: overhead setup



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Procedure: overhead setup



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Procedure: boiling setup

Top

Side



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Procedure: boiling setup



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Results



Crystallization of supersaturated Sodium Acetate solution from below. Three drops from an eyedropper. Two points of emanation. Shot at 1000 fps, f/8 90mm lens.



Crystallization of supersaturated Sodium Acetate solution from above. Crystals appear to be on top of the solution. Three drops from an eyedropper. Two points of emanation. Shot at 1000 fps, f/11, 90mm lens. Brightness increased 30% for clarity.

Results



Five frames of Sodium Acetate crystalizing from below. Each frame is .2 seconds apart, all five frames span .8 seconds. Shot at 1000 fps, f/8, 90mm lens.

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Results



Close up of solidified bismuth. Color variation from crystal formation is visible. Size is roughly ~3 mm across. f/16, 90mm lens

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Results



Solidified bismuth drop. Drop was melted in a spoon then poured onto a slide. The line across the middle of the image is the edge of the slide that the drop landed on. Many colors are visible along the various surfaces. f/36, 90mm lens.

Discussion: summary

Sodium Acetate:

- Crystals form radially outward from seed
- Crystals are long and skinny prisms
- Crystal formation pattern depends on how crystallization is triggered

Bismuth:

- Difficult to form nice looking bismuth crystals
- Colors won't show up on monochrome high-speed
- No crystal forming process static enough to record

Discussion: conclusions

- Sodium Acetate is feasible
 - Crystallizes consistently and on recordable time scales
- Bismuth isn't feasible
 - Crystallization isn't happening on time/physical scales we have the techniques to record

Discussion: moving forward

With sodium acetate:

- vary temperature/concentration, look at linear crystallization speed (see 1957 paper)
- look at types/processes of crystallization

Sodium Acetate







Sodium Acetate



Stable supersaturated solutions crystallize when seed is introduced

High Speed Crystallization of Sodium Acetate



Robin Deits, Sally Wolfe, Daniel Ron, Bayo Olatunji

03/01/12 •

High Speed Crystallization of Sodium Acetate



Robin Deits, Sally Wolfe, Daniel Ron, Bayo Olatunji

02/29/12 •

Temperature Regulation



Previous Work

• Dietz, Bruckner, and Hollingsworth (1957)



Previous Work

• Dietz, Bruckner, and Hollingsworth (1957)



How to Sound Like a Human (by forgetting your past) Alan Huang

CHALLENGE ACCEPTED







- Plot
- Conflict
- Structure
- Setting
- Theme
- Motif
- Style
- Imagery
- Symbolism
- Dialogue
- Dorapootivo



• Words



What comes after what?

Professor…

- 28.5% •McGonagall
- 8.9% •Trelawney
- •Umbridge 8.7%
- 5.9% •Dumbledore
- 5.8% •Lupin 4.6%
- •Snape
- •Flitwick

•S prout

from …

- 31.1% •the 6.1% •his
- •a
- •behind
- 2.5% •him
- •her

3.9%

3.3%

- •under
- •their

- 3.1% 2.9%
- 2.4%1.6%

1.6%

Professor Dumbledore knew something extremely odd stuff I'm going to feed the lake and Lies "Isn't it?"

Markov chain (the next thing you do depends only on the last few things you did) (in other words, you *forget* where you were before)



Longer phrases, up to a limit

 $Professor \ McGonagall \cdots$

- •was 5.5%
- •and 4.2%
- •had 3.4%
- •said 2.7%

McGonagall was… •now 6.1% •hurrying 6.1% •looking 6.1%

•right 6.1%

Professor Flitwick burst into flames and curled up, purring deeply. The common room emptied as people drifted off to bed. he went into the thick black trees.

Harry Potter and the Wide-Screen



This is actually useful!

Applications

- Polymer formation
- Cell behavior
- Statistical mechanics
- Financial modeling
- Societal evolution
- Data compression
- Music

Hidden Markov models

- Cryptography
- Error correction
- Speech recognition
- Computer vision
- Machine translation
- Fraud detection
- Pattern recognition

Google PageRank

- Internet contains N web pages
- Page contains k links
- Follow a link or type a URL
- Model as a Markov chain
- What pages do users end up on?
- Higher ad prices, earlier search results

Markov chains: *forgetful* models



PLAYING IN THE SANDBOX

David Benjamin

Process Isolation in Google Chrome



*** STOP: 0x00000019 (0x0000000,0xC00E0FF0,0xFFFFEFD4,0xC0000000) BAD_POOL_HEADER

CPUID: Genuine Intel 5.2.c irgl:1f SYSVER 0xf0000565

DII Base	DateStmp	- Name	DII Base	DateStmp -	- Name
80100000	32020070	- stockupl eve	80010000	31006052 -	- bal dll
99991999	310d06b4	- atani sus	888866888	31006074 -	2V2 TRADI 202
99206999	210J06bf		99201999	21012270 -	- Dick cuc
00200000	21006075		0020000	210010-7	
00201000	STecoc/a	- CLH332.313	00070000	Sieedoar	010000000
10098000	STecec/d	- Floppy.515	1 Cba8000	Slecbcal -	- Caron.sys
1 C 9 0 a 0 0 0	3lec6df7	- FS_Rec.SYS	10202000	316cecaa -	- Null.5Y5
fc864000	31ed868b	- KSecDD.SYS	fc9ca000	31ec6c78 -	- Beep.SYS
fc6d8000	31ec6c90	- 18042prt.sys	£c86c000	31ec6c97 -	- mouclass.sys
fc874000	31ec6c94	- kbdclass.sys	fc6f0000	31f50722 -	- VIDEOPORT.SYS
feffa000	31ec6c62	- mga_mil.sys	fc890000	31ec6c6d -	- vga.sys
fc708000	31ec6ccb	- Msfs.SYS	fc4b0000	31ec6cc7 -	- Npfs.SYS
fefbc000	31eed262	- NDIS.SYS	a0000000	31f954f7 -	- win32k.sys
fefa4000	31f91a51	- mga.dll	fec31000	31eedd07 -	 Fastfat.SYS
feb8c000	31ec6e6c	- TDI.SYS	feaf0000	31ed0754 -	- nbf.sys
feacf000	31f130a7	- tepip.sus	feab3000	31f50a65 -	- netbt.sys
fc550000	31601a30	- e159x.sus	fc560000	31f8f864 -	- afd.sus
fc718888	31ec6e7a	- nethios.sus	£c858000	31ec6c9h -	- Parport .sus
£c878888	31ec6c9h	- Pavallel SYS	fc954000	31ec6c9d -	- PanUdm SYS
fc5b0000	31ec6ch1	2V2 Iciwa2 -	£0340000	31£5003b -	- ndn sus
faa2b000	21£7.1ba		60912000	22021 abo -	ent cue
reaspood	orrearpa	hup.sgs	restatooo	Szosrabe	31.4.333
Adamage	duond due	Build [1201]			- Nama
fac22JOA	00142000	001/2000 001//000	000 00033133	220202	- VSaann SUS
001471-0	00143200	00143000 00144000	11011000 000	000001	
00147108	00144000	C000000 11011000	C0300000 000	140	- ntoskrn1.exe
8014/1dC	80122000	10003100 10300000	e133C4D4 e13	330440	- ntoskrni.exe
80147304	80302310	00000230 00000034	00000000 000		– ntoskini.exe

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Restart and set the recovery options in the system control panel or the /CRASHDEBUG system start option.


What if a webpage crashes?



What if a webpage crashes?



Tab isolation

Chrome brings the same isolation to webpages

One crashed tab doesn't kill the browser





- Crashes mean browser bugs
- Why not just write correct code?

5,000,000

Playing in the sandbox

- How tab isolation works
- Sandboxing
- Principle of least privilege





- Tasks organized into processes
- Typically one process per application
- Each process has private memory, state, etc.
- Independently scheduled and killed



Operating System

Single-process browsing









- Each tab in a separate renderer process
- Renderers coordinated by browser process
- Communicate by sending messages
- Renderer bugs are recoverable
 - Browser process crashes are still fatal

Playing in the sandbox

- How tab isolation works
- Sandboxing
- Principle of least privilege



Browser security

- Webpages are very restricted
 - Cannot read your English paper
 - Cannot write a virus
 - Cannot delete your files
- Visiting evil.example.com should be safe
- Browser enforces these restrictions
- What if there is a bug?

Enforcing security

- Sandbox each renderer
 - Can't access files
 - Internet
 - Other processes
- Only communicate with browser
 - Performs privileged actions on behalf of renderer
 - Including network access

Enforcing security



Enforcing security



Playing in the sandbox

- How tab isolation works
- Sandboxing
- Principle of least privilege



Principle of least privilege

Each component of your system should have only the rights it needs and no more

Limit damage when things go wrong
Fewest privileges to parts most likely to fail

Playing in the sandbox

- □ All complex software is buggy, so design for bugs
- Separate renderer process per tab
 - Improved stability
 - Security
- Restrict renderer rights to bare minimum
- Principle of least privilege
- Questions?

What do we mean by "innovation"? *Purpose driven creativity*

An abstract sculpture?



Creative.

An abstract sculpture that prevents the wind from spinning a building's revolving doors too fast?



Innovative.

Distributed Innovation There are a lot more smart people than geniuses

1 in a million?



Millions and billions.



New technologies enable collective genius

Culture and Persuasion



Finding Possible Proverbs



Inspirations For Helical Training

Alternate Reality Games

Large-scale exercises



What Skills Could Be Learned?



Sample Tasks





What would you learn if you walked through LA wearing 30 pounds of carrots, led a squirrel fishing party, challenged a dragon, asked for help inflating a large duck, and solicited strangers for a dinner invitation?







WOULD YOU INVITE TWO STRANGERS TO YOUR HOME FOR DINNER?



Spinner Assembly: Design





Not Shown: -Micro DC Motor -15 Tooth Brass Gear 15mm PD -Small Gear Mounting Collar -Flexible Coupling



Payload Mounting Plate Flexible Coupling Btaring Protusing Multi-Load Bearing Hollow Shaft 65 Tooth Steel Gear 65mm 2500 Tick Encoder Disk Insert Spacers Bus Cover Plate

Review

Optical Encoder Slip Ring Spacer Encoder Mounting Bracket

Slip Ring





Specifications

Max Angular Velocity: 1.7 Hz

Motor Controller: Pololu QiK TTL Interface

Slip Ring Channels: 12

Dimensions (Interface Space): 100mm x 100mm x 18mm



nstitute of



Massachusetts Institute of Technology

Spinner Assembly:







19 July 2011



Results:

Massachusette

Noticeable signal distortion appears at frequencies >1 KHz

Significant signal distortion appears at signal frequencies >1 MHz





1,000,000 bytes sent; 0 errors

Error probability <0.0000054 with 99% confidence

Conclusion:

Slip ring interface is suitable for serial data







Spinner Assembly redesign was successful

Meets volume, data transmission and pointing knowledge requirements

Next Step: Selection of Space Rated Components


The Photoelectric Effect



Explanations-Classical and Quantum

In 1887 Hertz observes a new phenomena

No classical explanation exists In 1905 Albert Einstein proposes a quantum explanation

 $E_{photon} = h\nu \quad KE_{max} = h\nu - \Phi$

Experimental Setup



Raw Data-Voltage vs. Current

275 Current, Voltage pairs collected over 5 different wavelengths



Finding the Cutoff Voltage



Graphical Results



Numerical Results

Method	Plank's Constant (Js)	Work Function (eV)
Linear Fit High Freq.	(1.07 ±0.37) · 10 − 34	0.03±0.175
X-Intercept High Freq.	(2.04±0.37)·10−34	-0.45±0.175
X-Intercept All Data	(6.72±1.40)-10-35	0.20±0.06
Actual	6.63.10-34	2.3

Experimental results differ from known values by at least an order of magnitude (several standard deviations)

Systematic Error

Potassium build up on the anode causes reverse photoelectric currents Insufficient Higher Voltage Data

Questions ?

The SPHERES IDE







The SPHERES Process

